

Hoolet: An OWL Reasoner with Support for Rules



Sean Bechhofer, Ian Horrocks

University of Manchester

<http://owl.man.ac.uk/hoolet>

Reasoning with OWL

- OWL DL has a “standard” first-order style semantics
- This allows us to use known results from Description Logic research to build reasoners for OWL
 - FaCT, RACER, Pellet
- However, the expressiveness of “full” OWL DL causes some problems
 - Currently no know **effective** algorithms in the presence of **cardinality, inverses and enumerations**
 - Reasoners such as FaCT and RACER “pretend” to handle one-of.
- Can we use alternative reasoning engines?

OWL and First Order Reasoning

- An alternative approach is to translate OWL DL into equivalent FOL axioms and then use a FO prover to provide inference
- Disadvantages
 - In general this compromises **decidability**, although a FO reasoner **may** be able to apply a complete strategy.
 - DL reasoners have been specifically **optimised** to handle DL style reasoning tasks. FO reasoners may require extra tuning to handle the tasks created.
- Advantages
 - Can handle **all** of OWL DL
 - Can be extended to deal with language **extensions** such as SWRL.

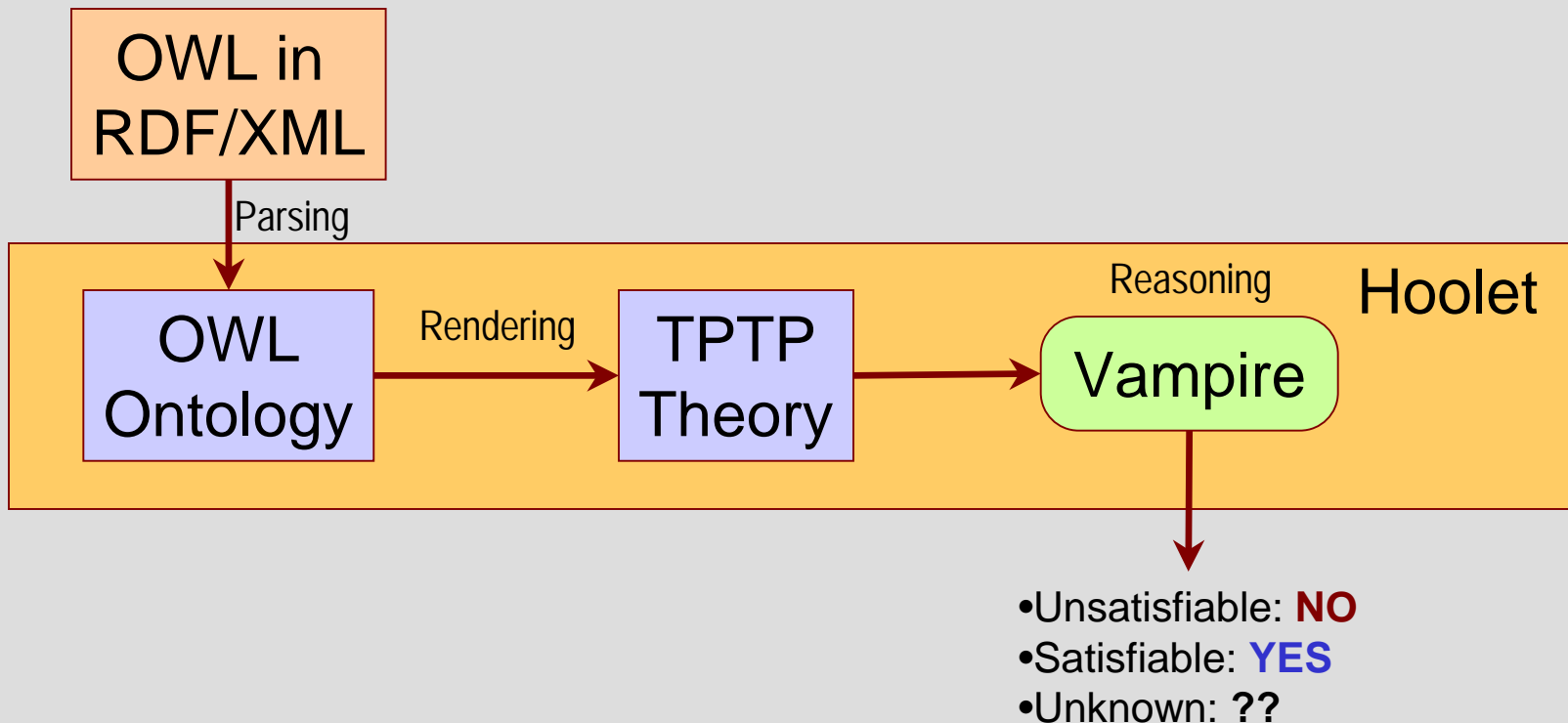
Hoolet

- A (prototype) OWL reasoner using a First Order prover.
- OWL ontology translated to equivalent axioms using the standard **TPTP** format.
- Axioms then passed to **Vampire** for satisfiability testing.
- Queries are translated to conjectures which are added to the theory.
- Hoolet may not be a very **effective** reasoner
 - This naive approach is not likely to scale well.
- However, it **does** provide a useful tool for use on small illustrative examples.
 - And may form **part** of an effective reasoning infrastructure

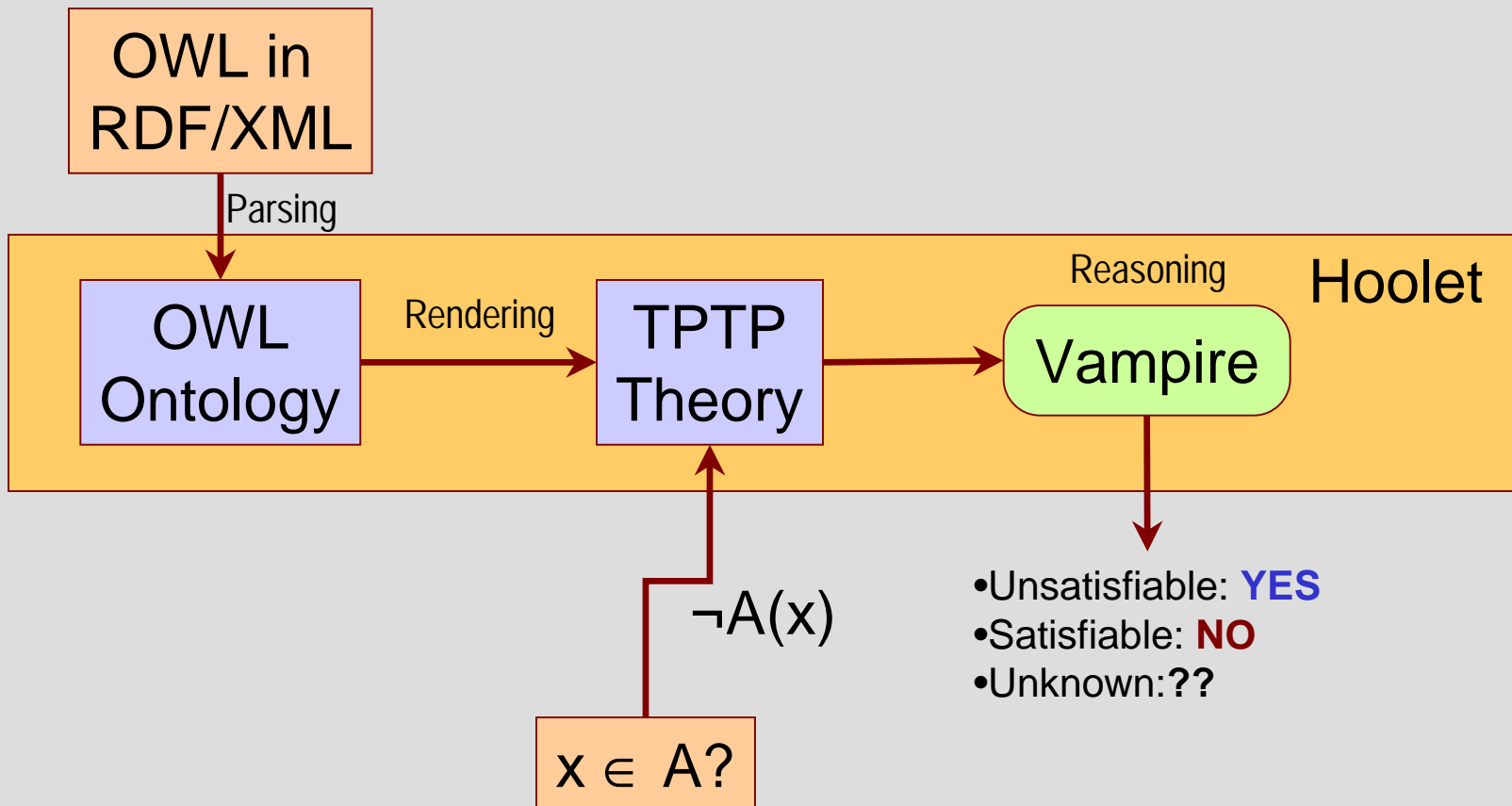
Example Translations

Class(B complete A)	$\exists x.A(x) \wedge B(x)$
SubClassOf(intersectionOf (A B) unionOf(C D))	$\exists x.(A(x) \wedge B(x)) \wedge (C(x) \vee D(x))$
Class(B partial restriction(p someValuesFrom A))	$\exists x.B(x) \wedge (\exists y.A(y) \wedge p(x,y))$
Class(A complete one-of(a b c))	$\exists x.A(x) \wedge (x=a \vee x=b \vee x=c)$

Satisfiability Testing



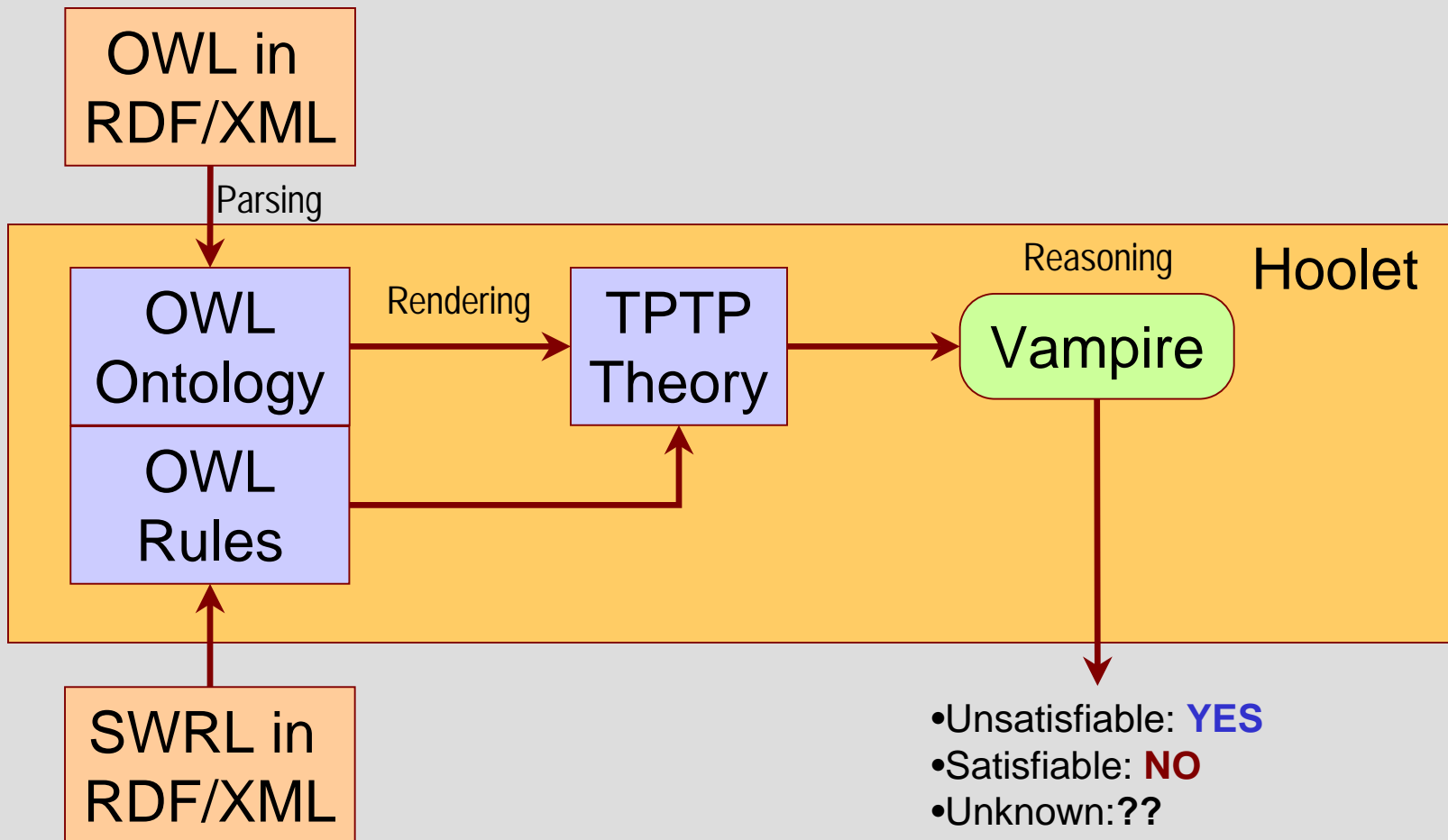
Query



Rules

- It is easy to extend Hoolet to handle **SWRL** rules.
- Each rule is simply translated to an **axiom** according to the semantics of the rules, with free variables universally quantified.
 $\text{hasParent}(\text{?x}, \text{?y}), \text{hasSibling}(\text{?y}, \text{?z}), \text{male}(\text{?z}) \rightarrow \text{hasUncle}(\text{?x}, \text{?z})$
translates to:
 $\forall x, y, z. \text{hasParent}(x, y) \wedge \text{hasSibling}(y, z) \wedge \text{male}(z) \rightarrow \text{hasUncle}(x, z)$
- Rules are then added to the theory.

Adding Rules



Hoolet Application

- Hoolet supplies a simple GUI for loading ontologies and rules
 - Uses **WonderWeb OWL API** for parsing and representation.
 - (Ab)uses **Vampire** prover for reasoning.
- Ontologies should be represented using OWL in RDF/XML
- Rules are represented using a (possibly idiosyncratic) RDF schema.
 - **Restrictions** on rule atoms: only classes allowed.
- Simple Queries:
 - satisfiability
 - subsumption
 - retrieval.
- Prototype from **<http://owl.man.ac.uk/hoolet>**